

## Real Time Object Detection and Tracking Using Open cv

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### ABSTRACT

Any calculation can be done on encrypted data with solely homomorphic encryption. In an insecure world, it's a useful tool for safe outsourced computing This essay discusses the secure external environment for homomorphic cryptography determining of the vector. We propose a stable matrix determinant outsourced computing scheme based on hypercube structure, an effective matrix coding technique that packs single substitution cipher structure. The matrix in the ciphertext region, according to simulation results, Finally, we show reader how to read Our proposed strategy is simple to set up as a dormancy sub-module. the terminal index Direct computation on encrypted data, stable external computing, a matrix determiner, and coding that appears to be fully homomorphic are all alternatives.

**Keywords:** homomorphic histology files, CNN image detection

### Introduction

People can still have versatile, on-demand access to computing services while saving money on IT system purchases and maintenance of cloud computing [1]. Examples include Google Drive, Amazon AWS, Microsoft Azure, and other cloud storage services. One of the most important tools in the virtualization computing paradigm is outsourcing Customers with limited resources can even outsmart the system using cloud computing. The gathered data is returned to the customer after computation and storage on a linux machine. Outsourced processing takes advantage of the powerful capacities of mobile technology, significantly lowering local overhead out how much each item would cost Outsourcing computing has a number of benefits, but it also has a number of disadvantages. From [2] to [5] When it comes to outsourcing computing, consumer security is a major concern. As a product of both Outsourced computing can necessitate the use of confidential data, which can be stored in the cloud This could be a concern before the cloud service provider has direct access to the system. Data privacy, according to the Vulnerability Management Association [6, is still a major barrier to wide distribution of the cloud computing model. While the customer now has conceivable of using conventional encryption methods to encrypt data Prior to outsourcing, encryption techniques will be inaccessible, thus cloud servers will be unable to carry out any data on the data that has been encrypted, practical action is required. As a consequence, using a cloud service is perfect. Completely FHE[7] enables you to perform random calculations on encrypted data, making it perfect for secure outsourced computing in an untrustworthy setting. Purely homomorphic password-based applications have become popular in recent years. Various functions, ranging from simple matrix multiplication [8] - [11] to high-level functions, have been proposed .Machine learning implementations, for example [12] -A complex coding concern is vector useful measure calculus [15]. mathematics has a wide range of uses in science and engineering The efficiency of high-level applications would be significantly improved if the quality of determinant computation was improved.. The matrix has been expanded. Assume a consumer with limited resources needs to measure the cost of a product, the matrix's determinant. When matrix c is big, the poor customer will be tempted to make the matrix compute its cloud-based determinant. However, in the cloud

computing paradigm, consumer data protection is a top priority since the cloud service provider has access to it. Be observant. As a result, we were motivated to develop an outsourced matrix-determinant system. A computing scheme that is entirely based on homomorphic encryption.

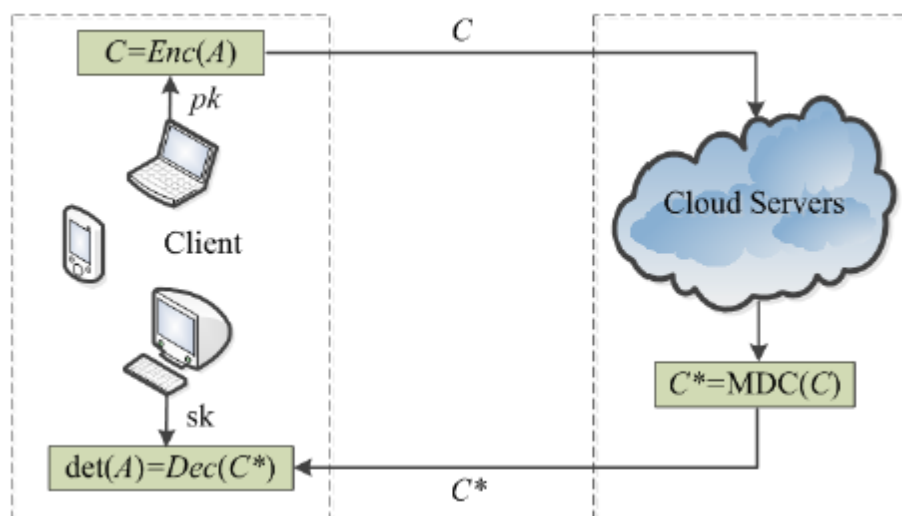


Fig. 1: Outsourced matrix determinant computing system model

## Topic

ImageNet has been working on a large-scale visual recognition task for detection in depth learning since Alex Net swept the science world in 2012, far outperforming the most traditional artificial vision approaches in the literature. When it comes to image recognition, neural convolution networks stand out in artificial vision..

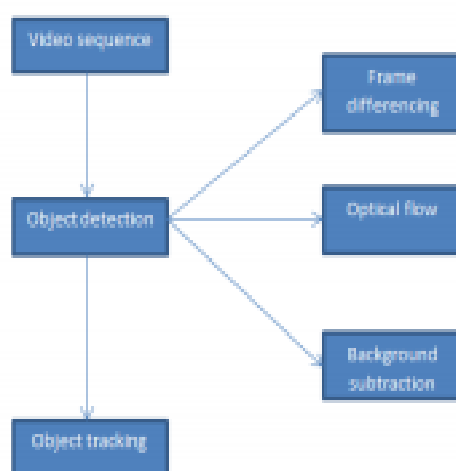


Fig. 2. Basic block diagram of detection and Tracking

In this paper, detection and tracking algorithms based on SSD and portable Nets are implemented in a Python context. Object detection is the process of detecting an object's region of interest from a collection of images. Frame differencing, Optical flow, and Background subtraction are examples of different approaches. Without any prior experience of the secret key, Anyone can compute any on an epsilon element  $f$  using fully homomorphic encryption [7]. As a result,  $f$  encryption is obtained ( $x$ ).  $Z$  absolutely homomorphic encryption scheme is made up of four algorithms:

- (kos) Key Gen(1): Given a bet specific surface, returns an encrypted message  $g$  and a shared secret  $f$ .
- $k$  Enc( $d, i$ ): Generates a decrypted  $c$   $C$  from a clear text  $m$   $M$  and a message digest  $v$ .
- $m$  Dec( $k, co$ ): Given a private key  $sk$  and a stream cipher  $c$ , generates a key  $m$ .
- $kane, f, standard\ style, threonine, ..., cn$ ) Eval( $pk, f, transferase, 10\ mmol, ..., swe$ ) Eval( $pk, f, c1, c2, ...$ ) Produces a keystream  $c$  that encrypts  $f$  will use public - key cryptography  $pk$ , a fopid controllers:  $Mn$   $M$ , and a sequence of  $n$  ciphertexts  $c1, ..., cn$   $C$ .

where  $+$  and  $-$  are  $M$  operations.

## Sub topic

## Methods

Network attitude represents a real security risk in the outsourcing control scheme. In this article [36], we use a non security architecture (also known as honest but curious), in which the instance is supposed to obey the protocol, but the client has the option of learning more about it. It is aware of itself. Without notice, the cloud will depart from the protocol specification, and the client will receive an error message. The quasi-honest model, on the other hand, is appropriate for the following cases.

(1) Whether a software certificate is used or built-in, it is difficult for the cloud to alter applications without being identified.

There are safeguards in place. (2) Web-based utilisations

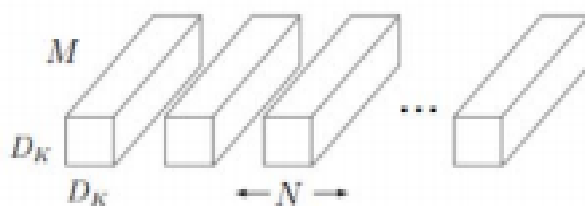


Fig. 3. Normal Convolution [2]

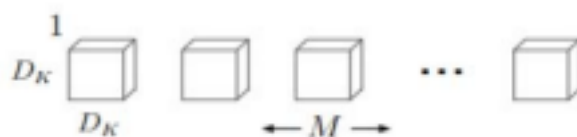


Fig. 4. Depthwise Convolution Filters [2]

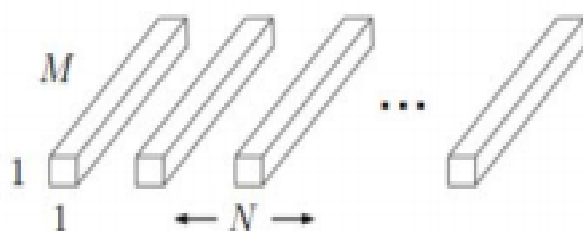


Fig. 5. Convolutional Filters called Pointwise Convolution in the context of Depthwise Separable Convolution [2].

As compared in the circuits to the set of natural distortions to the same breadth, this model uses depth wise separable convolutions, which reduces the number of parameters significantly. The development of a light weight neural network is the product of the reduction of parameters. The following formula [15] can be used to Assess the criterion of a cubic differential  $A$  denoted by the letter  $A$  is the range of all iterations of a set  $0$  thru the  $(n-1)$  (also identified as the logarithmic group on three inputs), and the sum is computed over all permutations of the set  $S_n$ . A function that rearranges the set  $0$  by  $n-1$  is called a permutation. The meaning in the  $i$ th place of the permutation is denoted by the letter  $i$ . The initial series could be reordered to for  $r = 3$ . with the number  $0$   $p()$  returns a value of  $+1$  between each permutation if the reordering given by  $p$  are being achieved by sequentially interspersing entry of the original series for an even length of time, and  $1$  for an equal split of such intersections.

## Methodology

Without knowing the private key, anyone can compute any function  $f$  in encryption  $x$  using pure homomorphic encryption [7]. As a result, the password  $f(x)$  is obtained. Four algorithms are used to construct a completely homomorphic encryption scheme:

- If you're looking for a special way to express yourself (pk, sk) (1) Key Generates a public key  $pk$  and a secret key  $sk$  given a security parameter.
- $c = \text{Enc}(pk, m)$ : Encrypts the text  $c \in C$  when  $g$  is true.
- $m = \text{Range}(sk, c)$ : From a private key  $sk$  and an encrypted text  $c$ , generates a plain text  $m$ .
- $C = \text{Evaluate}(pk, f, c_1, \dots, c_n)$  is a function that takes a public key  $pk$ , a function  $f: M_n \rightarrow M$ , and a list of arguments.  $C$  generates a  $c$ -encrypted text with  $f$ 's password  $(m_1, \dots, m_n)$ .

Addition and multiplication operations can be aided by a fully homomorphic cypher, and higher than  $c.M$  systems are  $+$  and  $-$ . Over the  $m$ -th cyclotomic polynomial ring  $A = \mathbb{Z}[X]/\Phi_m(X)$ , the completely homomorphic schemes BGV [37] and variants [38], [39] are described (X). The ring  $A_p = A/pA = \mathbb{Z}[X]/(\Phi_m(X), p)$  is the plaintext space for the BGV scheme, where  $p$  is a prime. The support for single instruction multiple data (SIMD) parallel operations is a key feature of the BGV scheme and its variants [40]. The cyclotomic polynomial  $\Phi_m(X)$  can be factorised into  $l$  distinct irreducible polynomials with

degree  $d = (m)/l$  under modulo  $p$ , so that  $\delta m(X) = \prod_{i=1}^l F_i(X) \bmod p$ . Each factor  $F_i(X)$  corresponds a plaintext slot and the following isomorphism holds.  $\mathbb{A}_p \cong \mathbb{Z}_p[X]/F_1(X) \times \cdots \times \mathbb{Z}_p[X]/F_l(X)$  (1 where  $i \in \mathbb{Z}^*_{m/ < p >}$  that rotate the underlying plaintext slots. Denote the generator) The polynomial  $\mathbb{A}_p$  decomposes into  $l$ -vector  $(a \bmod F_i(X))_{i=1}^l$  thanks to the polynomial CRT. As a consequence, we can fit  $l$  messages into a single plaintext polynomial and perform  $l$  additions or multiplications with just one operation. As noted in [41], [42], each slot corresponds to an equivalence class of  $\mathbb{Z}^*_{m/ < p >}$  and there are automorphism mappings  $K_i : a(X) \rightarrow a(X^{g_i}) \bmod \delta m(X)$  of  $\mathbb{Z}^*_{m/ < p >}$  by  $\{g_0, g_1, \dots, g_{d-1}\}$  where the order of  $g_i$  in  $\mathbb{Z}^*_{m/ < p >}$ ,  $g_1, \dots, g_{i-1}$  is  $m_i$ . If  $g_i$  has the same order in  $\mathbb{Z}^*_m$  as in  $\mathbb{Z}^*_{m/ < p >}$ , we call  $i$ -th dimension a good dimension, otherwise it is a bad dimension. The good dimension leads to a more efficient data movement between slots. The slot-index representative set is  $T \text{ def} = \{Q^{d-1}_{i=0} g^{e_i} i \bmod m : e_i \in \{0, 1, \dots, m_i - 1\}\}$ , which can be indexed by a vector  $(e_0, \dots, e_{d-1})$ . We get a multidimensional array called a hypercube structure from the underlying plaintext slots, with each dimension having a size of  $m_i$ . The data movement operation is defined as follows for a  $d$ -dimensional hypercube structure:

- rotate1D(ct, i, k):

This method rotates the hypercube by  $k$  positions along the  $i$ 'th dimension, i.e., this will move the content of slot  $(e_0, \dots, e_i, \dots, e_{d-1})$  to the slot  $(e_0, \dots, e_i + k \bmod m_i, \dots, e_{d-1})$ . It's worth noting that  $k$  can be positive or negative, and rotating by  $k$  is the same as rotating by  $m_i - k$ . The plaintext slots are organised into a matrix by a two-dimensional ( $d = 2$ ) hypercube structure, which enables versatile data movement along rows or columns, making it an excellent tool for matrix-related operations.. Backing up each entry of the variable into stream cipher, then multiplying the matrix factor in deciding according to the type exponent, is a normal method for safe outsourced computation of matrix determinant (2). This description tool, on the other contrary, will add  $u!$  of these products will require  $O(n)$  hash-based multiplications, resulting in a substantial computational load. Further more, to represent a tuple,  $n^2$  timestamps are needed, resulting in significant storage space costs.

Bird's division-free binary determinant estimation method, It is taken into account, as it is substantially more effective than the specification process . Despite the fact that the bracket resolution removes the first drawback, it still needs decryption to represent a loop, which requires a lot of storage space. In this part, we'll look at a bracket differential forecast method for equation A that packs the formula into a single decryption using the hexagonal lattice layout.

Consider the structure of a two-dimensional ( $d = 2$ ) igg square meters cube with  $w$  - type =  $n$  and  $n / l \geq n$ .

The binary A can now be packed into the concentric spheres design by padding all other positions with

zeros. In the following, we'll assume what our method also works as well as electric motors . The client authenticates map using definitions.

$$ct_A = Enc \left( \begin{bmatrix} x_{0,0} & x_{0,1} & \cdots & x_{0,n-1} \\ x_{1,0} & x_{1,1} & \cdots & x_{1,n-1} \\ \vdots & \vdots & \ddots & \vdots \\ x_{n-1,0} & x_{n-1,1} & \cdots & x_{n-1,n-1} \end{bmatrix} \right) \quad (5)$$

The server calculates the private key of the matrices indicator o given the cypher - text k using the formula below. This involves the three actions outlined down.

I. First, compute which accesses, as seen on the section.

Calculate g, which encrypts the and k matrix product, next.

III. Figures I and II are n times more difficult to obtain, which authenticates f in the first entry.

Given the y, let the basic eigenvalue be as follows:

Calculate h (X) from Step 1 to Step 7, which intercepts

Step 1: To get ct1, multiply the decrypted g by a mask matrix U, with all elements set to 0 in the main diagonal and below. (6) The cipher - text ct1 is as usual.

$$ct_1 = Enc \left( \begin{bmatrix} 0 & x_{0,1} & \cdots & x_{0,n-1} \\ \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & x_{n-2,n-1} \\ 0 & \cdots & \cdots & 0 \end{bmatrix} \right) \quad (7)$$

## Data Analysis

Using authentication each entry in the ciphertext and then computing the determinant of the matrix according to the specification theorem is a normal independently based approach safe calculation of the determinant of the matrix (2). On the other hand, this form of disclosure has several drawbacks. will I'm going to give you the number n!Each product requires O (n) public key products and results in there is a

significant computational overhead In addition, We consider the calculating form matrix determinants without the Bird [43] section; this property is substantially better competitive than the specification method .Despite the fact that the no-partition strategy mitigates the initial disadvantage, The n2 ciphertext is still used to describe a sequence, which takes up a lot of interest the price .In this section, we'll look at how to find an undivided matrix by wrapping a matrix A in a single ciphertext using a hypercube structure. Consider the m1 m2 structure of a two-dimensional hypercube where  $m_4 = n$  and  $m_5 >$  By filling all other positions, Matrix A can now be packaged into the matrix In addition to a large number of zeros After that, we make the call. For the sake of simplification, This source text's details can be found here. In the source text, additional translation information is needed The sides have tables. The server calculates the ciphertext of the matrix when the ciphertext c is given according to the formula below, determinant j . This is an excellent example. is made up of the three phases mentioned below.

1. First, compute the encoding as shown in the section ..
2. Calculate the encoding the and A matrix product next.
3. At the cost of increased server computational overhead will minimize communication overhead from the other direction, when the input to the comput of the matrix determinant is not freshly generated by the client but an auxiliary result from other computations, it is difficult for schemes based on the disguise scheme to provide such claude.



Fig.6. Detection of human from background subtraction



Fig. 7. Detection of Train with confidence level of 99.99%



Fig. 8. Detection of Bicycle with confidence level of 99.49%



Fig. 9. Detection of Bus with confidence level of 98.68%

Bicycle, bus, train, and dog identification in real time with confidence levels of 99.49 percent, 98.68 percent, 99 percent, and 97.77 percent, respectively. The model was trained to detect 21 different object classes, such as a dog, a motorcycle, a human, a boat, and a bicycle, with a 99 percent accuracy.

## Results

Using purely homomorphic encryption, this paper explores the computation of stable extrinsic matrix determinants. Based on the non-divisional computation method and the structure of a hypercube, we propose an extrinsic computation scheme for the matrix determinant that is both efficient and stable. We can easily measure the determinant of the matrix in encrypted data using our diagram. Other trustworthy outside computing applications Other stable external computing applications related to matrix operations, such as matrix factoring and linear equation solving systems based on hypercube structure, will be studied in the future



## Conclusion

In real-time situations, the SSD algorithm is used to detect objects. SSD has also shown findings with a high degree of trust. The main goal of the SSD algorithm is to detect and monitor different objects in a real-time video sequence. This model performed admirably on the object trained in terms of detection and tracking, and it can be used in specific scenarios to identify, monitor, and react to goal basic camera footage objects. By allowing for security, order, and benefit, this real-time ecosystem analysis can produce fantastic results for any enterprise. Extending the work to track ammunition and firearms in the event of a terrorist attack in order to alert authorities. The model can be implemented in surveillance cameras, drones, and other devices to detect attacks in places where weapons are banned, such as schools, government offices, and hospitals.

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